

# The PageRank Citation Ranking: Bring Order to the Web

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# Outline

- Introduction
- A ranking for every page on the Web
- Implementation
- Convergence Properties
- Personalized PageRank
- Conclusion

# Introduction

- Information Retrieval encounter many new challenges in the large and heterogeneous World Wide Web.
- Traditional IR:  
importance of academic papers  $\Rightarrow$  academic citation analysis.
- Web pages vs. academical publications:
  - Quality of academical publications: strictly controlled  
Web pages: free of quality control or publishing costs
  - $\hookrightarrow$  Huge number of pages can be created easily and inflating citation counts can be artificially manipulated
  - Academic papers: well defined units of work (quality & citation)  
Web pages vary on a wider scale than academic papers in quality, usage, citation and length

$\implies$  Relative importance of web pages: **PageRank**

# Link Structure of the Web

Link structure of the Web:

- pages = nodes & links = edges
- forward links = outedges
- backlinks = inedges
- A and B are Backlinks of C

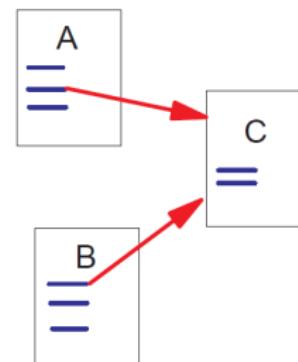


Figure: Link structure of page A, B and C

# Link Structure of the Web

Standard citation analysis for PageRank:

- a link from page A to B is a vote from A to B
- highly linked pages are more “important” than pages with few links
- backlinks from “important” pages are more significant than links from average pages

⇒ PageRank: how good an approximation to “importance” can be obtained just from the link structure of web



# First Version PageRank: Simplified Ranking Function

Summarize: a page can get a high rank if it has high sum of the ranks of its backlinks.

**First version of PageRank (simplified ranking function):**

$$R(u) = c \sum_{v \in B_u} \frac{R(v)}{N_v}$$

- $u$  is a web page
- $F_v$ : set of pages  $v$  points to
- $B_u$ : set of pages pointing to  $u$
- $N_v = |F_v|$
- $c$ : normalization factor

## First Version PageRank: Simplified Ranking Function

Assumptions of simplified ranking function:

- Every page must be reachable by some other page(s) in the web graph through link structure

Page ranking in the random surfer model:

more backlinks (from important pages)  $\Rightarrow$  more chance to be reached by random click

# First Version PageRank: Simplified Ranking Function

Another way to state simplified ranking function of PageRank:

- $A$ : a square matrix (columns and rows corresponding web pages)
- $A_{v,u}$ : link structure between page  $u$  and  $v$   
 $A_{v,u} = \frac{1}{N_u}$  if there is edge from  $u$  to  $v$   
 $A_{v,u} = 0$  otherwise
- $R$ : vector of PageRank over all pages

$$R = cAR$$

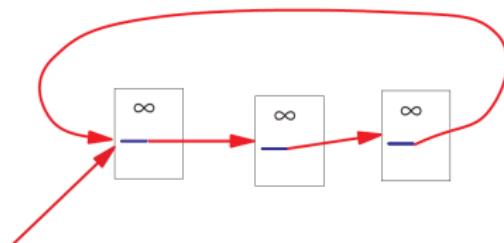
→  $R$  is an eigenvector of matrix  $A$  with eigenvalue  $c$ .

# Rank Sink

Simplified ranking function is constructed according to the intuitive basis: amount and importance of the backlinks

But: the web graph in real world is more complicated!

\* A special situation:



- Loop only with inedge (no outedges)
- Iteration of PageRank: Loop accumulate rank (no distribution further)
- Consequence: pages in loop with high ranking (ranks of pages outside the loop are zero)

⇒ **rank sink**

**Figure:** Loop which acts as a "rank sink"

# Rank Sink

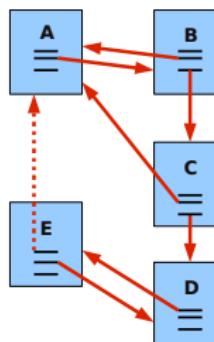


Figure: Loop accumulate rank (if without arrow from page E to A)

Loop (pages E and D) accumulate rank in calculation iteration

Consequently page E and D: PageRank 0.5 respectively

Other pages (A, B, C): zero

⇒ Distortion of the real importance ranking of pages

## Developed Version PageRank

### Definition

Let  $E$  be some vector over the Web Pages that corresponds to a source of rank. Then, the PageRank of a set of Web pages is an assignment,  $R'$ , to the Web pages which satisfies

$$R'(u) = c \sum_{v \in B_u} \frac{R'(v)}{N_v} + cE(u)$$

such that  $c$  is maximized and  $\|R'\|_1 = 1$  ( $\|R'\|_1 = 1$  denotes the  $L_1$  norm of  $R'$ ).

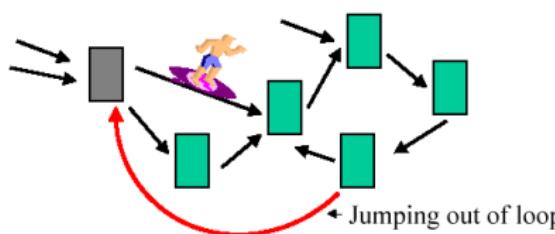
$E$ : Supplement the page rank for every page (so PageRanks not only rely on the link structure);

Defined uniform over all web pages (in most test) or user define

## $E$ vector in Random Surfer Model

Another intuitive basis of new Version in Random Surfer Model:  
Simplified version supposes that the random surfer will keep  
clicking on links as long as links are valid, but in the real world:  
if surfer get into a loop  $\Rightarrow$  jump out to visit another random page

- \* Which page will be randomly chosen to jump to based on the distribution in  $E$
- \*  $\|E\|_1 = 0.15$



# Dangling Links

**Dangling links:** links point to any page without outgoing links  
↪ "dead ends" with no further distribution of rank

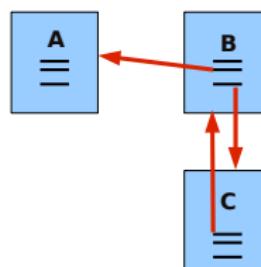


Figure: Dangling Links (page A is "dead end")

Solution: remove them from system until PageRanks for all of other pages are calculated, then put them back into system and recompute PageRanks.

# Computing PageRank

## Algorithm:

$R_0 \leftarrow S$  *# initialize vector over web pages*

loop:

$R_{i+1} \leftarrow AR_i$  *# new ranks sum of normalized backlink ranks*

$d \leftarrow \|R_i\|_1 - \|R_{i+1}\|_1$  *# compute normalizing factor*

$R_{i+1} \leftarrow R_{i+1} + dE$  *# add escape term*

$\delta \leftarrow \|R_{i+1} - R_i\|$  *# control parameter*

while  $\delta > \varepsilon$  *# stop when converged*

# Implementation

**Computing resources:** database with 24 million web pages  
references 75 million unique URLs

## Memory and disk:

- \* Weight vector: each URL 4 byte float (75 million URLs = 300MB)
- \* Weights from current time step stored in Memory, previous Weights accessed linearly on disk
- \* Link database Matrix A also in Disk (linear access)

## Implementation:

- ① Unique integer ID for each URL
- ② Sort and Remove dangling links
- ③ Rank initial assignment
- ④ Iteration until convergence
- ⑤ Add back dangling links and Re-compute
- ⑥ Took ca.5 hours

# Convergence Properties

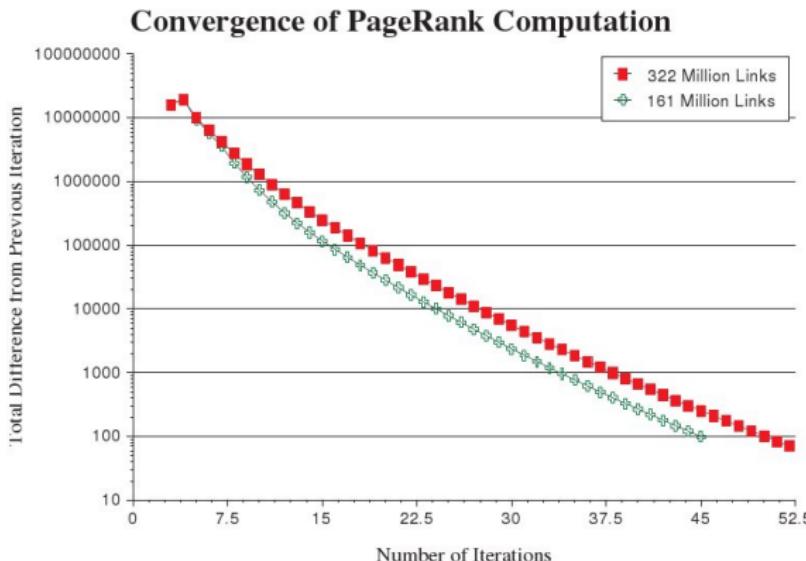


Figure: Rates of Convergence for Full Size and Half Size Link Databases;  
Scaling factor: linear in  $\log n$

# Convergence Properties

- PageRanks computation in logarithmic time
- Web Graph is rapidly mixing: an expander
- Properties of expander graphs can be utilized in future computation involving the Web graph

# Personalized PageRank

## $E$ Vector:

- Rank source to make up for the rank sinks
- The distribution of web pages that a random surfer periodically jumps to
- Customize page ranks

Test: the total weight  $E$  on a single web page

- \* Two home pages: Netscape & John McCarthy
- \* Result: two home pages get highest PageRank respectively, higher PageRanks belong to its immediate links
  - e.g. pages related to computer science at Stanford University get much higher PageRank in John McCarthy-rank than Netscape-rank

⇒ Setting vector  $E$  can generate personalized PageRanks

# Conclusion

## PageRank:

- Based just on Web's graph structure
- Citation analysis: backlinks (amount and importance)
- Developed PageRanks: rank source for a good approximation of “importance” ranking
- High quality search results

**Thank you for your attention!**